

MMT: A Foundation-Independent Approach to Declarative Languages

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My Background

▶ Areas

- ▶ theoretical foundations

logic, type theory, foundations of mathematics

- ▶ formal knowledge representation

specification, formalized mathematics, ontologies

- ▶ scalable applications

module systems, libraries, system integration

▶ Vision

- ▶ Develop a universal framework

for the formal representation of knowledge and its semantics,

- ▶ apply it to the safe and scalable integration of

math, logic, and computer science.

▶ Methods

- ▶ survey and abstract

understand fundamental concepts

- ▶ relate and transfer

unify different research areas

- ▶ long-term investment

identify stable ideas, do them right

- ▶ modularity and reuse

maximize sharing across languages, tools

Foundations

- ▶ Foundation = the most primitive formalism on which everything else is built
 - set theories, type theories, logics, category theory, ...
- ▶ We can fix the foundation once and for all — but which one?
- ▶ In math: usually implicit and arbitrary foundation
 - ▶ can be seen as avoiding subtle questions
 - ▶ but also as a strength: it's more general
- ▶ In CS: each system fixes its own foundational language
 - e.g., a variant of type theory or HOL

Fixed Foundations

- ▶ Fixing foundation the first step of most implementations
often foundation and implementation have the same name
- ▶ No two implementations for the exact same foundation
even reimplementations diverge quickly
- ▶ Negative effects
 - ▶ isolated, mutually incompatible systems
no sharing of results, e.g., between proof assistants
 - ▶ no large scale libraries
each system's library starts from scratch
 - ▶ no library archival
libraries die with the system
 - ▶ comparison of systems difficult
no common problem set
 - ▶ slow evolution
evaluating a new idea can take years

Overview

- ▶ Foundation-independent framework
 - ▶ avoid fixing foundation wherever possible
 - ▶ design and implement generically
 - ▶ permit instantiation with different foundations
- ▶ MMT language
 - ▶ prototypical formal declarative language
 - ▶ admits concise representations of most languages
 - ▶ continued development since 2006 (with Michael Kohlhase)
 - ▶ > 100 pages of publication
- ▶ MMT system
 - ▶ API and services
 - ▶ continued development since 2007 (with > 10 students)
 - ▶ > 30,000 lines of Scala code
 - ▶ ~ 10 papers on individual aspects

The Meta-Meta-Logical Approach

- ▶ Foundation-independence = last step in a progression of abstractions
- ▶ MMT governs the meta-meta-level
- ▶ In retrospect, I define MMT as Meta-Meta-Tool

Mathematics	Logic	Meta-logic	Foundation-Independence
			MMT
		logical framework	
	logic		
domain knowledge			

Foundation-Independence

- ▶ MMT arises by iterating the following steps
 1. Choose a typical problem
 2. Survey and analyze the existing solutions
 3. Differentiate between **foundation-specific** and **foundation-independent** concepts/problems/solutions
 4. Integrate the foundation-independent aspects into MMT
 5. Define interfaces to supply the foundation-specific aspects
- ▶ Separation of concerns between
 - ▶ foundation developers focus on logical core
 - ▶ service developers e.g., search
 - ▶ application developers e.g., IDE

yields rapid prototyping for logic systems

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- yields rapid prototyping for logic systems**
- ▶ But how much can really be done foundation-independently?
MMT shows: not everything, but a lot

Basic Concepts

Design principle

- ▶ few orthogonal concepts
- ▶ uniform representations of diverse languages

sweet spot in the expressivity-simplicity trade off

Concepts

- ▶ theory = named set of declarations
 - ▶ foundations, logics, type theories, classes, specifications, ...
- ▶ constant = named atomic declaration
 - ▶ function symbols, theorems, rules, ...
 - ▶ may have type, definition, notation
- ▶ term = unnamed complex entity, formed from constants
 - ▶ expressions, types, formulas, proofs, ...
- ▶ typing $\vdash_T s : t$ between terms relative to a theory
 - ▶ well-formedness, truth, consequence ...

Small Scale Example (1)

Logical frameworks in MMT

```

theory LF {
  type
  Pi      #  $\Pi V1 . 2$                                 name[: type][#notation]
  arrow   #  $1 \rightarrow 2$ 
  lambda  #  $\lambda V1 . 2$ 
  apply   #  $1\ 2$ 
}

```

Logics in MMT/LF

```

theory Logic : LF {
  prop : type
  ded   : prop  $\rightarrow$  type #  $\vdash 1$                                 judgments-as-types
}
theory FOL : LF {
  include Logic
  term      : type                                higher-order abstract syntax
  forall    : (term  $\rightarrow$  prop)  $\rightarrow$  prop #  $\forall V1 . 2$ 
}

```

Small Scale Example (2)

FOL from previous slide:

```

theory FOL : LF {
  include Logic
  term      : type
  forall    : (term → prop) → prop #  ∀ V1 . 2
}

```

Algebraic theories in MMT/LF/FOL:

```

theory Magma : FOL {
  comp : term → term → term # 1 ∘ 2
}
theory SemiGroup      : FOL {include Magma, ...}
theory CommutativeGroup : FOL {include SemiGroup, ...}
theory Ring : FOL {
  additive : CommutativeGroup
  multiplicative : Semigroup
  ...
}

```

Theories and Theory Morphisms

- ▶ Theories
 - ▶ uniform representation of
 - ▶ foundations e.g., logical frameworks, set theories, ...
 - ▶ logics, type theories
 - ▶ domain theories e.g., algebra, arithmetic, ...
 - ▶ little theories: state every result in smallest possible theory
maximizes reuse
- ▶ Theory morphisms
 - ▶ uniform representation of
 - ▶ extension e.g., Monoid \rightarrow Group
 - ▶ inheritance e.g., superclass \rightarrow subclass
 - ▶ semantics e.g., FOL \rightarrow ZFC
 - ▶ models e.g., Nat: Monoid \rightarrow ZFC
 - ▶ implementation e.g., specification \rightarrow programming language
 - ▶ translation e.g., typed to untyped FOL
 - ▶ functors e.g., output \rightarrow input interface
 - ▶ homomorphic translation of expressions
 - ▶ preserve typing (and thus truth)

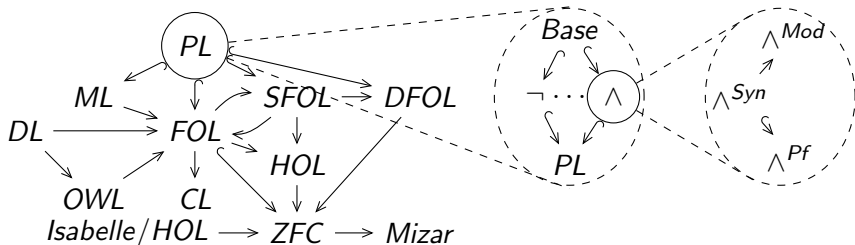
Large Scale Example: The LATIN Atlas

- ▶ DFG project 2009-2012 (with DFKI Bremen and Jacobs Univ.)
- ▶ Highly modular network of little logic formalizations
 - ▶ separate theory for each
 - ▶ connective/quantifier
 - ▶ type operator
 - ▶ controversial axioms e.g., excluded middle, choice, ...
 - ▶ base type
 - ▶ reference catalog of standardized logics
 - ▶ documentation platform
- ▶ Written in MMT/LF
- ▶ 4 years, with ~ 10 students, ~ 1000 modules

Logic Diagrams in LATIN

An example fragment of the LATIN logic diagram

- ▶ nodes: MMT/LF theories
- ▶ edges: MMT/LF theory morphisms



- ▶ each node is root for library of that logic
- ▶ each edge yields library translation functor

library *integration* very difficult though

Current State

- ▶ Little theories including
 - ▶ propositional, common, modal, description, linear logic, unsorted/sorted/dependently-sorted first-order logic, CASL, higher-order logic
 - ▶ λ -calculi (λ -cube), product types, union types, ...
 - ▶ ZFC set theory, Mizar's set theory, Isabelle/HOL
 - ▶ category theory
- ▶ Little morphisms including
 - ▶ relativization of quantifiers from sorted first-order, modal, and description logics to unsorted first-order logic
 - ▶ negative translation from classical to intuitionistic logic
 - ▶ translation from type theory to set theory
 - ▶ translations between ZFC, Mizar, Isabelle/HOL
 - ▶ Curry-Howard correspondence between logic, type theory, and category theory

Logical Results Obtained at the MMT Level

- ▶ Module system
modularity transparent to foundation developer
- ▶ Concrete/abstract syntax
notation-based parsing/presentation
- ▶ Interpreted symbols, literals
external model/implementation reflected into MMT
- ▶ Type reconstruction
foundation plugin supplies only core rules
- ▶ Simplification
rule-based, integrated with type reconstruction
- ▶ Theorem proving? very recent, ongoing
- ▶ Code generation? most likely

Abstract Syntax

Key ideas

- ▶ no predefined constants
- ▶ very general term constructor
- ▶ $c(\Gamma; \vec{E})$ binds variables and takes arguments
 - ▶ non-binding operators: Γ empty e.g., `apply(\cdot ; f, a)` for LF's $f a$
 - ▶ typical binders: Γ and \vec{E} have length 1
e.g., `lambda($x:A$; t)` for LF's $\lambda x:A.t$

contexts	$\Gamma ::= (x[: E][= E])^*$
terms	$E ::=$
constants	c
variables	x
complex terms	$ c(\Gamma; E^*)$

MMT implements foundation-independent data structures for theories and terms

Concrete Syntax

- ▶ Notations of c introduce connect concrete syntax for $c(\Gamma; \vec{E})$

e.g., for LF

production	MMT declaration	abstract syntax
$E ::=$		
type	type #	type
$\Pi x : E_1 . E_2$	Pi # $\Pi V1 . 2$	Pi($x : E_1 ; E_2$)
$E_1 \rightarrow E_2$	arrow # $1 \rightarrow 2$	arrow($\cdot ; E_1, E_2$)
$\lambda x : E_1 . E_2$	lambda # $\lambda V1 . 2$	lambda($x : E_1 ; E_2$)
$E_1 E_2$	apply # $1 2$	apply($\cdot ; E_1, E_2$)

MMT implements foundation-independent parser and printers

Inference System

For any theory Σ :

$\vdash_{\Sigma} \Gamma$	Γ is a valid context
$\Gamma \vdash_{\Sigma} t : A$	t has type A
$\Gamma \vdash_{\Sigma} E = E'$	E and E' are equal
$\Gamma \vdash_{\Sigma} - : A$	A is inhabitable

MMT define some foundation-independent rules

- ▶ congruence rules for equality
- ▶ contexts

$$\frac{}{\vdash_{\Sigma} \cdot} \quad \frac{\vdash_{\Sigma} \Gamma \quad [\Gamma \vdash_{\Sigma} - : A] \quad [\Gamma \vdash_{\Sigma} t : A]}{\vdash_{\Sigma} \Gamma, c[: A][= t]}$$

- ▶ rules for **atomic terms**, e.g.

$$\frac{x : A \text{ in } \Gamma}{\Gamma \vdash_{\Sigma} x : A} \quad \frac{x = t \text{ in } \Gamma}{\Gamma \vdash_{\Sigma} x = t}$$

Foundation-specific rules for **complex terms** are

- ▶ declared in theories
- ▶ implemented by plugins

Inference System: Implementation

MMT implements foundation-independent parts of type checker

- ▶ foundation-independent rules
- ▶ lookup in theories, context
- ▶ simplification, definition expansion
- ▶ error reporting

Foundation-specific rules supplied by plugins

- ▶ ~ 8 abstract rules, e.g.,
 - ▶ infer type
 - ▶ check term at given type
 - ▶ check equality of two given terms
 - ▶ simplify a term
- ▶ each rule can recurse into other judgements
- ▶ plugins provide concrete instances
- ▶ Example LF: ~ 10 rules for LF, ~ 10 lines of code each

Inference System: Type Reconstruction

Type Reconstruction

- ▶ Judgment with unknown meta-variables
 - ▶ implicit arguments, type parameters
 - ▶ omitted types of bound variables
- ▶ Goal: prove judgment and solve meta-variables
- ▶ Much harder than type checking **requires delaying constraints**

MMT implements foundation-independent type reconstruction

- ▶ transparent to foundations
- ▶ (almost) no extra cost for foundation developer
one additional rule for LF

Application-Independence

- ▶ Practical logic-related systems often application-**specific**
 - ▶ fixed functionality for fixed foundation
 - often: read-eval-print design
 - ▶ many applications shallow, decay quickly
- ▶ MMT approach: application-**independence**
 1. focus on API for MMT language
 - data structures and high-level interfaces
 2. encapsulate advanced functionality in reusable services
 - e.g., theorem proving, search, ...
 3. allow plugin interfaces for customization
 - foundation-specific rules, parsers, change listeners, ...
 4. build lightweight applications on top IDE, library manager, ...

Knowledge Management Results at the MMT Levels

- ▶ Change management recheck only if affected
- ▶ Project management indexing, building
- ▶ Extensible export infrastructure
Scala, SVG graphs, LaTeX, HTML, ...
- ▶ Search, querying substitution-tree and relational index
- ▶ Browser interactive web browser
- ▶ Editing IDE-like graphical interface

IDE

- ▶ Inspired by programming language IDEs
- ▶ Components
 - ▶ jEdit text editor (in Java): graphical interface
 - ▶ MMT API (in Scala)
 - ▶ jEdit plugin to tie them together

only ~ 1000 lines of glue code
- ▶ Features
 - ▶ outline view
 - ▶ error list
 - ▶ display of inferred information
 - ▶ type inference of subterms
 - ▶ hyperlinks: jump to definition
 - ▶ search interface
 - ▶ context-sensitive auto-completion: show identifiers that

IDE: Example View

The screenshot shows the jEdit IDE interface with the following components:

- Title Bar:** jEdit - C:\other\oaff\test\source\examples\pl.mmt
- Menu Bar:** File Edit Search Markers Folding View Utilities Macros Plugins Help
- Left Panel (Sidekick):**
 - Tree view for file: C:\other\oaff\test\so...
 - pl.mmt
 - theory PL
 - prop
 - ded
 - and
 - impl
 - equiv
 - type
 - definition
 - lambda
 - x
 - prop
 - y
 - prop
 - and
 - impl
 - y
 - ded

- Main Editor:**

```

1 namespace http://cds.omdoc.org/examples
2 theory PL : http://cds.omdoc.org/urtheories?LF =
3   prop : type
4   ded  : prop → type
5   and  : prop → prop → prop
6   impl : prop → prop → prop
7   equiv : prop → prop → prop
8   = [x,y] (x ⇒ y) ∧ ded
9

```
- Status Bar:** 1 error, 0 warnings
- Bottom Panel (Error List):**

C:\other\oaff\test\source\examples\pl.mmt (1 error, 0 warnings)

- 8: invalid object: http://cds.omdoc.org/examples?PL?equiv?definition: ded
 - argument must have domain type
 - http://cds.omdoc.org/examples?PL; x:prop, y:prop |- ded : prop
 - http://cds.omdoc.org/examples?PL; x:prop, y:prop |- prop→type = prop
- Bottom Bar:** 8,30 (mmt,sidekick,UTF-8)Smr oWV 27/3 Mb 4 error(s)19:50

An Interactive Library Browser

- ▶ MMT content presented as HTML5+MathML pages
- ▶ Dynamic page updates via Ajax
- ▶ MMT used through HTTP interface with JavaScript wrapper
- ▶ Features
 - ▶ interactive display e.g., inferred types, redundant brackets
 - ▶ smart navigation via MMT ontology
 - ▶ can be synchronized with jEdit
 - ▶ dynamic computation of content
 - ▶ e.g., definition lookup, type inference
 - ▶ graph view: theory diagram as SVG

Browser: Example View

The MMT Web Server

[Graph View](#)
[Search](#)
[Administration](#)
[Help](#)

code.google.com / p / hol-light / source / browse / trunk ? bool

Style: html5

- hollight
- ├ arith.omdoc
- ├ bool.omdoc
- ├ calc_int.omdoc
- ├ calc_num.omdoc
- ├ calc_rat.omdoc
- ├ cart.omdoc
- ├ class.omdoc
- ├ define.omdoc
- ├ ind_defs.omdoc
- ├ ind_types.omdoc
- ├ int.omdoc
- ├ iterate.omdoc
- ├ lists.omdoc
- ├ nums.omdoc
- ├ pair.omdoc
- ├ real.omdoc
- ├ realarith.omdoc
- ├ relax.omdoc
- ├ sets.omdoc

bool

T [show/hide type](#) [show/hide onedim-notation](#) [show/hide tags](#) [show/hide metadata](#)

T_DEF [show/hide type](#) [show/hide tags](#) [show/hide metadata](#)

TRUTH [show/hide type](#) [show/hide definition](#) [show/hide tags](#) [show/hide metadata](#)

^ [show/hide type](#) [show/hide onedim-notation](#) [show/hide tags](#) [show/hide metadata](#)

AND_DEF [show/hide type](#) [show/hide tags](#) [show/hide metadata](#)

=> [show/hide type](#) [show/hide onedim-notation](#) [show/hide tags](#) [show/hide metadata](#)

IMP_DEF [show/hide type](#) [show/hide tags](#) [show/hide metadata](#)

! [show/hide type](#) [show/hide onedim-notation](#) [show/hide tags](#) [show/hide metadata](#)

type $\{A : \text{holtype}\} (A \Rightarrow \text{bool}) \Rightarrow \text{bool}$

onedim-notation $\forall x: _ . a$ (precedence 0)

FORALL_DEF [show/hide type](#) [show/hide tags](#) [show/hide metadata](#)

type $\{A : \text{holtype}\} \vdash (! A) = \lambda P:A \Rightarrow \text{bool} . P = \lambda x:A . T$

Browser Features: 2-dimensional Notations

REAL_POW_DIV

show/hide type

show/hide definition

type $\vdash \forall x:\text{real} . \forall y:\text{real} . \forall n:\text{num} . \left(\frac{x}{y}\right)^n = \frac{x^n}{y^n}$

Browser Features: Proof Trees

The MMT Web Server
[Graph View](#) [Administration](#) [Help](#)

Style: html5 cds.omdoc.org / courses / 2013 / ACS1 / exercise_10.mmt ? Problem3

acs1_2013

- exercise_10.omdoc
 - Problem2
 - Problem3**
 - Problem4
- example
- latin
- lmfdb
- mathscheme
- mml
- openmath
- test
- tptp
- urtheories

theory Problem3 **meta** LF

include : <http://cds.omdoc.org/examples?FOLEQNatDed>

circ : term → term → term

e : term

R : ⊢ ∀ x x ∘ e ≐ x

C : ⊢ ∀ x ∀ y x ∘ y ≐ y ∘ x

L : ⊢ ∀ x e ∘ x ≐ x

$$= \left[\begin{array}{c} \frac{\frac{\frac{C\ e}{\vdash \forall y x \circ y \doteq y \circ e} \text{forallE } x}{\vdash e \circ x \doteq x \circ e} \text{forallE} \quad \frac{R\ x}{\vdash x \circ e \doteq x} \text{forallE}}{\vdash e \circ x \doteq x} \end{array} \right]$$

⊢ ∀ x e ∘ x ≐ x

- reconstructed types >
- implicit arguments >
- redundant brackets >**
- infer type
- simplify
- fold

- show
- hide

Enter an object over theory: <http://cds.omdoc.org/courses/2013>

[x] x ∘ e

analyze simplify

[x] x ∘ e

[x:term] term

Browser Features: Type Inference

The screenshot shows a theorem prover interface with several theorem definitions and a context menu for type inference.

The definitions are:

- FORALL_DEF** `show/hide type` `show/hide tags` `show/hide metadata`
`type {A:holtype} ⊢ (!A) = λP:A ⇒ bool. P = λx:A. T`
- `? show/hide type`
- EXISTS_DEF** `show`
- `✓ show/hide type`
- OR_DEF** `show/hide`
- F** `show/hide type`
`type bool`

The context menu for the first definition includes the following options:

- reconstructed types >
- implicit arguments >
- redundant brackets >
- infer type** (highlighted with a mouse cursor)
- simplify
- fold

The **infer type** option is expanded, showing the inferred type: `(A ⇒ bool) ⇒ bool`.

Browser Features: Parsing

Enter an object over theory:

analyze simplify

result: $[x] \forall y. \exists z. y =_{\text{num}} x + z$

inferred type: $\{x: \text{num}\} \text{bool}$

Example Service: Search

Enter Java regular expressions to filter based on the URI of a declaration

Namespace

Theory

Name

Enter an expression over theory

Use \$x,y,z:query to enter unification variables.

Search

type of **MOD_EQ**

$\vdash \forall m:\text{num} . \forall n:\text{num} . \forall p:\text{num} . \forall q:\text{num} . m = n + q * p \implies m \text{ MOD } p = n \text{ MOD } p$

type of **MOD_MULT_ADD**

$\vdash \forall m:\text{num} . \forall n:\text{num} . \forall p:\text{num} . (m * n + p) \text{ MOD } n = p \text{ MOD } n$

L^AT_EX Integration

- ▶ MMT declarations spliced into L^AT_EX documents
shared MMT-L^AT_EX knowledge space
- ▶ L^AT_EX macros for MMT-HTTP interface
- ▶ Semantic processing of formulas
 - ▶ parsing
 - ▶ type checking
 - ▶ semantic enrichment: cross-references, tooltips
- ▶ Design not L^AT_EX-specific
e.g., integration with word processors possible

L^AT_EX Integration: Example

Inferred arguments are inserted during compilation:

- ▶ upper part: L^AT_EX source for the item on associativity
 - ▶ lower part: pdf after compiling with L^AT_EX-MMT
 - ▶ type argument M of equality symbol is inferred and added by MMT
-

```
\begin{mmtscope}
  For all \mmtvar{x}{in M}, \mmtvar{y}{in M}, \mmtvar{z}{in M}
  it holds that !(x * y) * z = x * (y * z)!
\end{mmtscope}
```

A *monoid* is a tuple (M, \circ, e) where

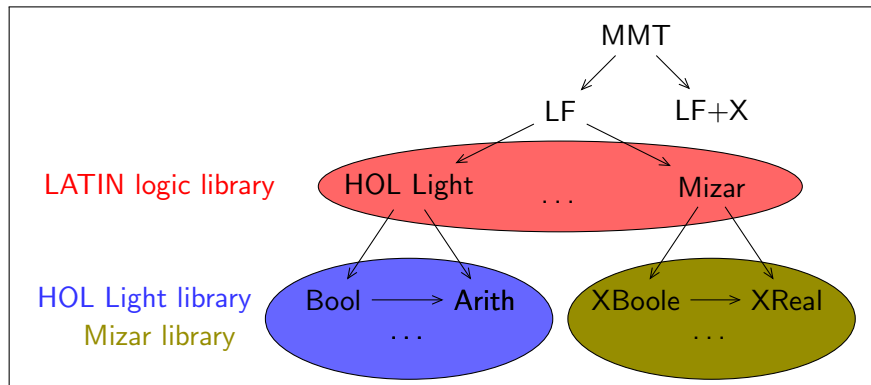
- M is a sort, called the universe.
- \circ is a binary function on M .
- e is a distinguished element of M , the unit.

such that the following axioms hold:

- For all x, y, z it holds that $(x \circ y) \circ z =_M x \circ (y \circ z)$
 - For all x it holds that $x \circ e =_M x$ and $e \circ x =_M x$.
-

Current Work: Library Integration

- ▶ Open Archive of Formalizations open PhD position!
Michael Kohlhase and myself, 2014-2017
- ▶ Goal: archival, comparison, integration of formal libraries
Mizar, HOL systems, IMPS, Coq/Matita, PVS, ...
- ▶ Big, overlapping libraries – that are mutually incompatible



Goal: Universal Library Infrastructure

- ▶ MMT as representation language
- ▶ Repository backend: MathHub
 - ▶ based on GitLab – open-source analog of GitHub server
 - ▶ GitLab instance hosted at Jacobs University
 - ▶ free registration of accounts, creation of repositories
- ▶ Generic library management
 - ▶ browser
 - ▶ inter-library navigation
 - ▶ search
 - ▶ change management

Goal: Towards Library Integration

- ▶ Refactor exports to introduce modularity
- ▶ 2 options
 - ▶ systematically during export
e.g., one theory for every HOL type definition
 - ▶ heuristic or interactive MMT-based refactoring
- ▶ Collect correspondences between concepts in different libraries
heuristically or interactively
- ▶ Relate isomorphic theories across languages
- ▶ Use partial morphisms to translate libraries

Conclusion

- ▶ MMT: general framework for declarative languages
 - ▶ Foundation-independent representation language
 - ▶ Application-independent implementation
- ▶ Easy to instantiate with specific foundations
 - ▶ rapid prototyping logic systems
- ▶ Multiple deep foundation-independent results
 - ▶ logical: parsing, type reconstruction, module system, ...
 - ▶ knowledge management: search, browser, IDE, ...
- ▶ MMT quite mature now, ready for larger applications
 - ▶ about to break even
- ▶ Interesting for
 - ▶ new, changing foundations
 - ▶ generic applications/services
 - ▶ system integration/combination

Further Resources

Web sites

- ▶ MMT: <https://svn.kwarc.info/repos/MMT/doc/html/index.html>
- ▶ OAF web server (experimental): <http://mathhub.info:8080/>

Selected publications all available from <https://kwarc.info/people/frabe>

- ▶ the primary paper on the MMT language (I&C 2013, with M. Kohlhase):
A Scalable Module System
- ▶ a more recent paper on the MMT approach to logic (JLC 2014):
How to Identify, Translate, and Combine Logics?
- ▶ Foundations in LATIN: (MSCS 2011, with M. Ianu)
Formalizing Foundations of Mathematics
- ▶ Modular logics in LATIN (TCS 2011, with F. Horozal):
Representing Model Theory in a Type-Theoretical Logical Framework
- ▶ Mizar in OAF (JAR 2013, with M. Ianu, M. Kohlhase, J. Urban):
The Mizar Mathematical Library in OMDoc: Translation and Applications
- ▶ HOL Light in OAF: (CICM 2014, with C. Kaliszyk)
Towards Knowledge Management for HOL Light